## Final Exam: December 17 Name:

You will have 3 hours for this exam, although you should not need that much. This exam is closed-book and closed-note. Please take some time to check your work. If you need extra space, write on the back. There are a total of 40 points on this exam.

- 1. (4 points) Consider the list of numbers
  - $41 \quad 17 \quad 23 \quad 4 \quad 36 \quad 19 \quad 12 \quad 28 \quad 5$
  - (a) After each of the first three passes of Insertion Sort, the list will be

Pass 1:	17	41	23	4	36	19	12	28	5
Pass 2:	17	23	41	4	36	19	12	28	5
Pass 3:	17	23	41	4	36	19	12	28	5

What will the list be after the next pass?

(b) After each of the first three passes of Selection Sort, the list will be

Pass 1:	5	17	23	4	36	19	12	28	41
Pass 2:	5	17	23	4	28	19	12	36	41
Pass 3:	5	17	23	4	12	19	28	36	41

What will the list be after the next pass?

2. (4 points) Given the ordered list of numbers

$$14 \quad 16 \quad 17 \quad 22 \quad 27 \quad 31 \quad 42$$

(a) What sequence of numbers will be examined in performing a sequential search for the target 28?

(b) What sequence of numbers will be examined in performing a binary search for the target 28?

3. (6 points) Complete the following C++ function which computes the binomial coefficient C(n,k). Use the following facts to write the function:

```
• C(n,0) = C(n,n) = 1;
  • If 0 < k < n, then C(n,k) = C(n-1,k-1) + C(n-1,k).
int C(int n, int k)
// Precondition: 0 <= k <= n \,
// Postcondition: Returns binomial coefficient "n choose k" --
     number of ways to choose a set of k out of n items
11
{
 if (
                                      ) return 1;
 return C(
                              ) + C(
                                                       );
                    ,
                                             ,
}
```

4. (4 points) Show the recursion tree (that is, the "box trace" as desribed in the text) for evaluating C(4, 2).

5. (10 points) Here is another (less obvious) C++ implementation of the binomial coefficient function:

```
int C2(int n, int k)
{
    int result = 1;
    for (int i = 1; i <= k; i++) {
        result = result * (n+1-i) / i;
    }
    return result;
}</pre>
```

You should check that C(4, 2) and C2(4, 2) give the same result.

- (a) What is the big-O running time to evaluate C2(n, n/2), expressed in terms of n?
- (b) Suppose that C2(100, 50) takes 50 ms to run on a particular machine; estimate how long C2(102, 51) will take:

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(c) An approximate recurrence relation for the running time of the original, recursive version on C(n, n/2) is:

$$\begin{bmatrix} T(0) = 1\\ T(n) = 2 \cdot T(n-1) \end{bmatrix}$$

Solve this recurrence:

- (d) Suppose that C(100, 50) takes 50 ms to run on a particular machine, using the recursive code; estimate how long C(102, 51) will take:
- (e) Which implementation is likely to run faster for large values of n and k, your recursive solution or this iterative version? Why?
- 6. (3 points) What is the output of the following sequence of operations if x is declared as a stack<char>?

```
x.push('h'); x.push('e'); x.push('l');
cout << x.top(); x.pop();
cout << x.top(); x.pop();
x.push('l'); x.push('o');
cout << x.top(); x.pop();
cout << x.top(); x.pop();
cout << x.top(); x.pop();</pre>
```

What would the output be from the same sequence if x were instead declared as a queue<char>, where now x.push(c) inserts c into the queue ("enqueues" it), x.top() returns the character at the front of the queue (in the C++ standard library, this operation is actually called x.front()), and x.pop() removes the character from the front ("dequeues" it).

- 7. (3 points) Consider the stack ADT as discussed in class: it only provides the operations push, pop, top, and empty (which returns true if the stack is empty), plus the usual constructor and destructor. Suppose we add the following three operations:
  - int size() return the number of items on the stack
  - StackItemType retrieve(int i) return the item at position i from the top of the stack
  - void replace(int i, StackItemType x) replace the item at position i with x

Give one benefit and one disadvantage to using this modified stack ADT.

8. (6 points) What is the output from the following code?

```
struct node {
  int item;
  node *next;
};
int a, b;
int *p, *q;
node *x, *y;
a = 25;
b = 17;
p = &a;
q = new int;
*p = a + b;
*q = a + b;
cout << "a = " << a << ", b = " << b << endl;
cout << "*p = " << *p << ", *q = " << *q << endl;
p = q;
q = 0;
cout << "*p = " << *p << endl;
x = 0;
y = new node;
y \rightarrow item = 5;
y \rightarrow next = x;
x = new node;
x \rightarrow item = 3;
x \rightarrow next = y;
for (y = x; y != 0; y = y - > next) {
  cout << y->item << endl;</pre>
}
```

Now write a few more lines to "clean up" after the above code by deleting all of the dynamically allocated storage: