Practice Exam 2

Solutions

This is a collection of relevant problems I have given on previous exams; it does not correspond exactly to a one-hour test. This exam is closed-book and closed-note. Please take some time to check your work. If you need extra space, write on the back. You must show your work to receive any partial credit.

1. Given the sets

$$A = \{a, i, u\}$$

 $B = \{i, o, u\},$

list the elements of each of the following sets:

(a)
$$A \cup B$$

(b) $A \cap B$

(c) A - B

(d) $(A - B) \cup (B - A)$

(e) $A \times B$

$$\{(a,i), (a,o), (a,u), (i,i), (i,o), (i,u), (u,i), (u,o), (u,u)\}$$

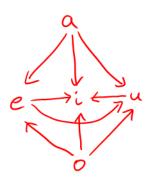
(f) P(A)

$$\{\phi, \{a\}, \{i\}, \{u\}, \{a,i\}, \{a,u\}, \{i,u\}, \{a,i,u\}\}\}$$

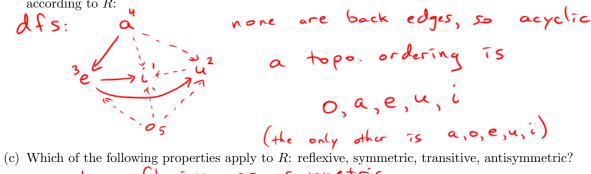
2. Let A be the set $\{a, e, i, o, u\}$, and consider the relation R on A whose graph is given by the following adjacency matrix:

(Recall that the convention is that the cell at row x, column y is 1 if x R y.)

(a) Draw the graph of R:



(b) Either identify a cycle in the graph of R, or give a topological ordering of the elements of A according to R:



3. Given the set $S = \{A, B, C, D, E, F, G\}$, we may represent any subset of S by its characteristic vector, which will have seven bits. For example, the set $\{A, B, D\}$ is represented by the bit vector 1101000. Consider these named subsets of S:

$$I = \{C, E, G\} \qquad iii = \{E, G, B\}$$
 On the bit vector representation for each of the following:

(a)
$$I \cup iii = \{B, C, E, G\}$$

$$0(1010)$$
or 0(0010)

(b)
$$I \cap iii = \{E, G\}$$

0010101

AND 0100101

(c)
$$I - V = \{C, E\}$$
AND $\frac{0.010101}{0.010100} = V$ (complement)

(d)
$$(I-V)\cup(V-I) = \{C,E\}\cup\{B,D\}$$

$$= \{B,C,D,E\}$$

$$= \{B,C,D,E\}$$
(e) $(I\cup V)-(I\cap V) = \{B,C,D,E,C\}-\{G\}$

$$= \{B,C,D,E\}$$
Same as (d)
$$= \{B,C,D,E\}$$

4. For the same set S, how many elements are in the powerset, P(S)? What is the set of bit vectors corresponding to the elements of P(S) (give a simple description)?

$$|S| = 7$$
, So $|P(s)| = 2^7 = 128$
the bit vectors are the 128 7-bit binary
numbers from 0000000 (o) to 1111111 (127).

5. We observed in class that the logical implication operator, \rightarrow , behaves like a transitive relation. Consider the set \mathcal{E} of all logical expressions. If E and F are elements of \mathcal{E} , then we will define the relation \Rightarrow on \mathcal{E} by saying that $E \Rightarrow F$ exactly when the expression $E \to F$ is a tautology. For example, $(p+q) \Rightarrow (q+p)$. However, $(p+q) \not\Rightarrow pq$, because when p is true and q is false, the condition p+q is true but the conclusion pq is not.

(a) Doe
$$pq \Rightarrow (p+q)$$
 hold?
 $yes - whenever pq$ is true, it is also the case that $p+q$ is true.

(b) Is \Rightarrow reflexive? Why or why not?

(c) Is \Rightarrow symmetric? Why or why not?

no — by the above,
$$Pq \Rightarrow P+q$$
 but $p+q \Rightarrow Pq$

(d) Is \Rightarrow antisymmetric? Why or why not?

no — for example, both
$$p+q \Rightarrow q+p$$

and $q+p \Rightarrow p+q$,

but $p+q \neq q+p$ (as expressions; they are however logically equivalent)

s \Rightarrow transitive? Why or why not?

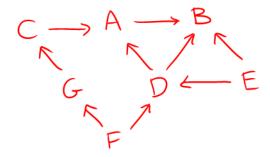
(e) Is \Rightarrow transitive? Why or why not?

yes — if
$$E \rightarrow F$$
 and $F \rightarrow G$ are tautologics,
then so is $E \rightarrow G$, for all E, F, G .

6. Here is an adjacency list representation of a directed graph:

Node	Successors
A	В
В	(none)
\mathbf{C}	A
D	A, B
\mathbf{E}	B, D
F	D, G
G	C

(a) Draw the graph.



(b) If the graph is acyclic, give a topological ordering of its nodes; otherwise, identify a cycle. If starting at $f: C_3^3 \to A_2^2 \to B_2^2$ no back edges, so acyclic topo. ordering: (one of several) E,F,D,G,C,A,B

(c) What is the longest path in the graph that never revisits a node?

(by inspection)
$$F \rightarrow G \rightarrow C \rightarrow A \rightarrow B$$

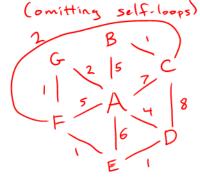
(d) Define the distance between two vertices as the length of the shortest path between them, or ∞ if there is no such path. What is the greatest non-infinite distance in this graph?



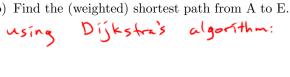
7. Here is an adjacency matrix representation of an undirected weighted graph (a weight of ∞ means there is no edge between those vertices):

	A	В	\mathbf{C}	D	\mathbf{E}	\mathbf{F}	\mathbf{G}
A	0	5	7	4	6	5	2
В	5	0	1	∞	∞	∞	∞
$^{\mathrm{C}}$	7	1	0	8	∞	2	∞
D	4	∞	8	0	1	∞	∞
\mathbf{E}	6	∞	∞	1	0	1	∞
\mathbf{F}	5	∞	2	∞	1	0	1
G	2	$\begin{array}{c} 5 \\ 0 \\ 1 \\ \infty \\ \infty \\ \infty \\ \infty \end{array}$	∞	∞	∞	1	0

(a) Draw the graph.



(b) Find the (weighted) shortest path from A to E.



- (c) Would the answer to the previous question change if the weight of the edge between B and C were changed to -1? Why or why not?

no change — the shortest path from A to either B or C is 5 units, so the -1 between B&C couldn't make a path shorter than 4 (in fact, the best path (d) Find a minimum spanning tree for the graph. using that edge is A-B-C-F-E this uses Prim's algorithm (like Dijkstra w/ different priority):

